

## Lecture 13. Support Vector Machine II

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## 1 The Primal Problem

Recall from the last lecture that, we are interested in the problems that take the form of

$$\begin{aligned} \min_{\mathbf{x}} \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & \mathbf{g}(\mathbf{x}) \leq 0, \\ & \mathbf{h}(\mathbf{x}) = 0, \\ & \mathbf{x} \in X. \end{aligned} \tag{1}$$

We denote the *feasible set* of (1) by

$$D_0 = \{\mathbf{x} : \mathbf{g}(\mathbf{x}) \leq 0, \mathbf{h}(\mathbf{x}) = 0, \mathbf{x} \in X\}. \tag{2}$$

Each element in  $D_0$  is called a *feasible solution*. The *optimal function value* is

$$f^* = \inf_{\mathbf{x} \in D_0} f(\mathbf{x}). \tag{3}$$

**Assumption 1. Feasibility and Boundedness** *The feasible set is nonempty and the objective function is bounded from below, that is,*

$$-\infty < f^* = \inf_{\mathbf{x} \in D_0} f(\mathbf{x}) < \infty.$$

## 2 The Lagrangian Dual Problem

### 2.1 Weak duality

Recall from the last lecture that, for any  $\lambda \geq 0$ , we have

$$q(\lambda, \mu) \leq f^*.$$

This immediately leads to the result as follows.

**Theorem 1. Weak Duality Theorem** *We define the *dual optimal value* by*

$$q^* = \sup_{\lambda \geq 0, \mu} q(\lambda, \mu). \tag{4}$$

*Then, we have*

$$q^* \leq f^*. \tag{5}$$

The optimization problem in (4) is the so-called *Lagrangian dual problem*. As we have shown that the dual function  $q$  is concave, the Lagrangian dual problem is indeed equivalent to a *convex optimization problem* (why?).

Theorem 1 implies that, the dual optimal value is a lower bound of the optimal function value  $f^*$ . The difference between  $f^*$  and  $q^*$  is the so-called duality gap.

**Definition 1.** *Duality gap* is defined by

$$f^* - q^*.$$

**Remark 1.** Duality gap is a commonly used termination condition for a set of optimization algorithms.

In terms of the duality gap, we naturally have a few questions to ask.

**Question 1.** When is the duality gap zero, i.e.,  $q^* = f^*$ ?

**Question 2.** Suppose that the duality gap is zero, can we find **a set of**  $(\lambda^*, \mu^*)$  with  $\lambda^* \geq 0$  such that

$$q^* = q(\lambda^*, \mu^*) = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) = f^*?$$

**Question 3.** Suppose that the duality gap is zero, and there exists  $(\lambda^*, \mu^*)$  with  $\lambda^* \geq 0$  such that

$$q^* = q(\lambda^*, \mu^*) = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) = f^*.$$

Then, if  $\hat{\mathbf{x}}$  **minimizes**  $L(\mathbf{x}, \lambda^*, \mu^*)$ , that is,

$$\hat{\mathbf{x}} \in \underset{\mathbf{x} \in X}{\operatorname{argmin}} L(\mathbf{x}, \lambda^*, \mu^*), \quad (6)$$

can we say that,  $\hat{\mathbf{x}}$  is one of the optimal solutions to the primal problem, i.e.,

$$\hat{\mathbf{x}} \in \underset{\mathbf{x} \in D_0}{\operatorname{argmin}} f(\mathbf{x})?$$

**Question 4.** Suppose that the duality gap is zero, and there exists  $(\lambda^*, \mu^*)$  with  $\lambda^* \geq 0$  such that

$$q^* = q(\lambda^*, \mu^*) = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) = f^*.$$

Then, if  $\hat{\mathbf{x}}$  minimizes  $L(\mathbf{x}, \lambda^*, \mu^*)$ , that is,

$$\hat{\mathbf{x}} \in \underset{\mathbf{x} \in X}{\operatorname{argmin}} L(\mathbf{x}, \lambda^*, \mu^*), \quad (7)$$

and  $\hat{\mathbf{x}}$  **is feasible**, i.e.,  $\hat{\mathbf{x}} \in D_0$ , can we say that,  $\hat{\mathbf{x}}$  is one of the optimal solutions to the primal problem, i.e.,

$$\hat{\mathbf{x}} \in \underset{\mathbf{x} \in D_0}{\operatorname{argmin}} f(\mathbf{x})?$$

All of the subsequent discussions are trying to answer the above questions. The major tool is the geometric treatment we introduced last lecture, that is, we treat the Lagrangian as a linear function over the space where the set  $S = (g_1(\mathbf{x}), \dots, g_m(\mathbf{x}), h_1(\mathbf{x}), \dots, h_p(\mathbf{x}), f(\mathbf{x}))$  lies in.

**Remark 2.** The major motivation for introducing the Lagrangian is to transforming a **constrained** optimization problem with the feasible set  $D_0$  to an **(almost) unconstrained** optimization problem with feasible set  $X$ , while the optimal function value remains the same.

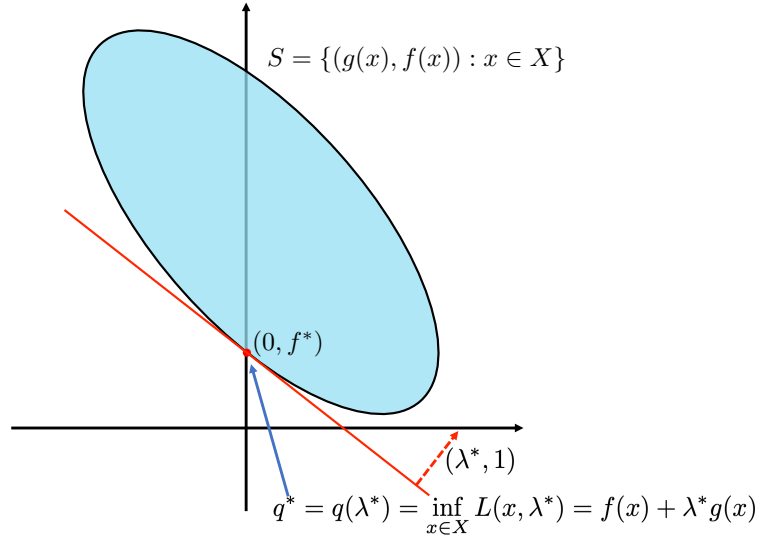


Figure 1: Illustration of the geometric multipliers.

## 2.2 The Geometric Multipliers

In view of Figure 1, the equality  $q^* = f^*$  holds implies that, we may be able to (but not necessarily, see [Remark 5](#)) find a hyperplane with the normal vector  $(\lambda^*, 1)$  that supports the set  $S$  from below intercepts the vertical axis at the level  $f^*$ . In this case, we can see that the duality gap is zero. This motivates the concept **geometric multipliers** as follows.

**Definition 2.** A vector  $(\lambda^*, \mu^*) = (\lambda_1^*, \dots, \lambda_m^*, \mu_1^*, \dots, \mu_p^*)$  is said to be a geometric multiplier vector (or simply geometric multiplier) for the primal problem if

$$\lambda_i^* \geq 0, i = 1, \dots, m,$$

and

$$f^* = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*). \quad (8)$$

**Remark 3.** Notice that, Eq. (8) is a requirement of the geometric multiplier instead of a definition of  $f^*$ . Recall that,

$$f^* = \inf_{\mathbf{x} \in D_0} f(\mathbf{x}).$$

**Remark 4.** The RHS of Eq. (8) is indeed  $q(\lambda^*, \mu^*)$ . Therefore, the existence of a geometric multiplier  $(\lambda^*, \mu^*)$  implies that we can find a feasible solution  $(\lambda^*, \mu^*)$  of the dual problem such that  $f^* = q(\lambda^*, \mu^*)$ .

**The existence of geometric multipliers** indeed implies that there is **no duality gap**. We formalize this result by the proposition as follows.

**Proposition 1.** Suppose that  $(\lambda^*, \mu^*)$  is a geometric multiplier vector of the primal problem. Then, we have the following hold.

1.  $q^* = q(\lambda^*, \mu^*)$ , that is,  $(\lambda^*, \mu^*)$  is one of the dual optimal solutions to the Lagrangian dual problem (4);

2. the duality gap is zero, i.e.,  $f^* = q^*$ .

*Proof.* Recall that, the Lagrangian dual function is defined by

$$q(\lambda, \mu) = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda, \mu).$$

Thus, the right hand side of Eq. (8) is indeed  $q(\lambda^*, \mu^*)$ , and we can write the condition in Eq. (8) as

$$f^* = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) = q(\lambda^*, \mu^*). \quad (9)$$

By further noting the weak duality property in (5) and the condition  $\lambda \geq 0$  in Definition 2, we can conclude that

$$q^* = q(\lambda^*, \mu^*), \quad (10)$$

that is, the geometric multiplier  $(\lambda^*, \mu^*)$  is one of the dual optimal solutions to the Lagrangian dual problem (4). Moreover, combining (9) and (10) immediately leads to  $f^* = q^*$ , which completes the proof.  $\square$

**Remark 5.** If we can find a geometric multiplier, then there is no duality gap. However, the converse is not true. That is, if there is no duality gap, we may not be able to find a geometric multiplier. They may not even exist at all.

**Example 1.** Consider an optimization problem as follows.

$$\begin{aligned} \min f(x) &= x \\ \text{s.t. } g(x) &= x^2 \leq 0, \\ x &\in X = \mathbb{R}. \end{aligned}$$

### 2.3 The Complementary Slackness

If a geometric multiplier  $(\lambda^*, \mu^*)$  is known, we hope that  $\hat{\mathbf{x}}$  that minimizes the Lagrangian  $L(\mathbf{x}, \lambda^*, \mu^*)$  over  $\mathbf{x} \in X$  is one of the optimal solutions to the primal problem as well. However, the vector  $\hat{\mathbf{x}} \in \operatorname{argmin}_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*)$  may not even be in the feasible set  $D_0$ .

**Example 2.** Consider an optimization problem as follows.

$$\begin{aligned} \min f(x) &= \begin{cases} e^x, & x \leq 0, \\ 1 - x, & x \in [0, 1], \\ 0, & x > 1. \end{cases} \\ \text{s.t. } g(x) &= x \leq 0. \end{aligned}$$

We can see that, the geometric multiplier  $\lambda^*$  is 0, and the corresponding Lagrangian is

$$L(x, \lambda^*) = f(x).$$

Thus,

$$\operatorname{argmin}_{x \in \mathbb{R}} L(x, \lambda^*) = \{x : x \geq 1\}.$$

Clearly, none of the points that minimizes  $L(x, \lambda^*)$  is feasible regarding the primal problem.

What if  $\hat{\mathbf{x}} \in \mathbf{argmin}_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*)$  is a feasible solution to the primal problem? Can we conclude that such a  $\hat{\mathbf{x}}$  is an optimal solution to the primal problem? The answer is still no.

**Example 3.** Consider an optimization problem as follows.

$$\begin{aligned} \min f(x) &= \begin{cases} -x, & x \leq 0, \\ 0, & x > 0. \end{cases} \\ \text{s.t. } g(x) &= x \leq 0. \end{aligned}$$

We can see that, the geometric multiplier  $\lambda^*$  is not unique and it can take any value from  $[0, 1]$ . Let us consider the case in which  $\lambda^* = 1$  and the corresponding Lagrangian is

$$L(x, \lambda^*) = f(x) + g(x) = \begin{cases} 0, & x \leq 0, \\ x, & x > 0. \end{cases}$$

Thus,

$$\mathbf{argmin}_{x \in \mathbb{R}} L(x, \lambda^*) = \{x : x \leq 0\}.$$

However, it is easy to see that only  $x^* = 0$  is the optimal solution to the problem.

**Remark 6.** Notice that, Example 3 also provides us an example that the geometric multiplier may not be unique. Indeed, for Example 3, the geometric multiplier is  $\lambda^* \in [0, 1]$ .

Thus, we need extra conditions to find the desirable optimal solutions from the set in (7), which is the so-called **complementary slackness**.

**Proposition 2.** Let  $(\lambda^*, \mu^*)$  be a geometric multiplier. Then,  $\mathbf{x}^*$  is a global minimum of the primal problem if and only if

$$\mathbf{x}^* \text{ is feasible,} \tag{11}$$

$$\mathbf{x}^* \in \mathbf{argmin}_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*), \tag{12}$$

$$\lambda_i^* g_i(\mathbf{x}^*) = 0, i = 1, \dots, m. \tag{13}$$

*Proof.*

1. ( $\Rightarrow$ ) Suppose that  $\mathbf{x}^*$  is a global minimum of the primal problem. Then,  $\mathbf{x}^*$  must be feasible, and thus

$$f(\mathbf{x}^*) \geq f(\mathbf{x}^*) + \sum_{i=1}^m \lambda_i g_i(\mathbf{x}^*) + \sum_{i=1}^p \mu_i h_i(\mathbf{x}^*) = L(\mathbf{x}^*, \lambda^*, \mu^*) \geq \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) = f^*.$$

The definition of  $f^*$  leads to  $f^* = f(\mathbf{x}^*)$ , which implies that the above inequality is an equality. Thus,

$$f(\mathbf{x}^*) = L(\mathbf{x}^*, \lambda^*, \mu^*) = f^* = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*).$$

This leads to (12) and

$$f(\mathbf{x}^*) = L(\mathbf{x}^*, \lambda^*, \mu^*) = f(\mathbf{x}^*) + \langle \lambda^*, \mathbf{g}(\mathbf{x}^*) \rangle + \langle \mu^*, \mathbf{h}(\mathbf{x}^*) \rangle.$$

As  $\mathbf{x}^*$  is feasible, that is,  $\mathbf{g}(\mathbf{x}^*) \leq 0$  and  $\mathbf{h}(\mathbf{x}^*) = 0$ , we have Eq. (13).

2. ( $\Leftarrow$ ) Suppose that  $\mathbf{x}^*$  is feasible and (12) and (13) hold.

In view of (12) and the fact that  $(\lambda^*, \mu^*)$  is the geometric multiplier, we have

$$L(\mathbf{x}^*, \lambda^*, \mu^*) = f^* = \inf_{\mathbf{x} \in D_0} f(\mathbf{x}).$$

Moreover, the feasibility of  $\mathbf{x}^*$  and (13) imply that

$$L(\mathbf{x}^*, \lambda^*, \mu^*) = f(\mathbf{x}^*) + \sum_{i=1}^m \lambda_i g_i(\mathbf{x}^*) + \sum_{i=1}^p \mu_i h_i(\mathbf{x}^*) = f(\mathbf{x}^*).$$

Combining the above two equations leads to

$$f(\mathbf{x}^*) = f^* = \inf_{\mathbf{x} \in D_0} f(\mathbf{x}),$$

which implies that  $\mathbf{x}^*$  is a global minimum of the primal problem.

The proof is complete. □

**Remark 7. Complementary slackness** in (13) implies that

$$\begin{aligned} \lambda_i^* > 0 &\Rightarrow g_i(\mathbf{x}^*) = 0, \\ g_i(\mathbf{x}^*) < 0 &\Rightarrow \lambda_i^* = 0. \end{aligned}$$

Complementary slackness is frequently used in characterizing the optimal solutions.

## 2.4 Primal and dual optimal solutions

**Theorem 2. Optimality Conditions** (The KKT Conditions) *A pair  $\mathbf{x}^*$  and  $(\lambda^*, \mu^*)$  is an optimal solution and geometric multiplier pair if and only if*

$$\mathbf{x}^* \in X, \mathbf{g}(\mathbf{x}^*) \leq 0, \mathbf{h}(\mathbf{x}^*) = 0, \quad (\text{Primal Feasibility}), \quad (14)$$

$$\lambda^* \geq 0, \quad (\text{Dual Feasibility}), \quad (15)$$

$$\mathbf{x}^* \in \underset{\mathbf{x} \in X}{\operatorname{argmin}} L(\mathbf{x}, \lambda^*, \mu^*), \quad (\text{Lagrangian Optimality}), \quad (16)$$

$$\lambda_i^* g_i(\mathbf{x}^*) = 0, i = 1, \dots, m, \quad (\text{Complementary Slackness}). \quad (17)$$

*Proof.*

1.  $\Rightarrow$  Suppose that  $\mathbf{x}^*$  and  $(\lambda^*, \mu^*)$  is an optimal solution and geometric multiplier pair. Then, the primal feasibility and dual feasibility hold.

Moreover,

$$f(\mathbf{x}^*) = f^* = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) \leq L(\mathbf{x}^*, \lambda^*, \mu^*) \leq f(\mathbf{x}^*),$$

which implies the Lagrangian optimality and the complementary slackness.

2.  $\Leftarrow$  Suppose that the conditions in (14) to (17) hold. Then

$$f(\mathbf{x}^*) = L(\mathbf{x}^*, \lambda^*, \mu^*) = \min_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) \leq \inf_{\mathbf{x} \in D_0} L(\mathbf{x}, \lambda^*, \mu^*) \leq \inf_{\mathbf{x} \in D_0} f(\mathbf{x}) \leq f(\mathbf{x}^*),$$

which implies that  $\mathbf{x}^*$  is the optimal solution and  $(\lambda^*, \mu^*)$  is the geometric multiplier.

The proof is complete.  $\square$

**Proposition 3. Saddle Point Theorem** (Optional) *A pair  $\mathbf{x}^*$  and  $(\lambda^*, \mu^*)$  is an optimal solution-geometric multiplier pair if and only if  $\mathbf{x}^* \in X$ ,  $\lambda^* \geq 0$ , and  $(\mathbf{x}^*, \lambda^*, \mu^*)$  is a saddle point of the Lagrangian, in the sense that*

$$L(\mathbf{x}^*, \lambda, \mu) \leq L(\mathbf{x}^*, \lambda^*, \mu^*) \leq L(\mathbf{x}, \lambda^*, \mu^*), \forall \mathbf{x} \in X, \lambda \geq 0, \mu \in \mathbb{R}^p. \quad (18)$$

*Proof.*

1.  $\Rightarrow$  As the pair  $\mathbf{x}^*$  and  $(\lambda^*, \mu^*)$  is an optimal solution-geometric multiplier pair, we have (14) to (17) hold. Clearly, we can see that  $\mathbf{x}^* \in X$ ,  $\lambda^* \geq 0$ , and the Lagrangian optimality in (16) implies that

$$L(\mathbf{x}^*, \lambda^*, \mu^*) \leq L(\mathbf{x}, \lambda^*, \mu^*), \forall \mathbf{x} \in X.$$

Moreover, in view of the definition of geometric multiplier, we have

$$f^* = \inf_{\mathbf{x} \in X} L(\mathbf{x}, \lambda^*, \mu^*) = L(\mathbf{x}^*, \lambda^*, \mu^*).$$

Thus, combining the feasibility of  $\mathbf{x}^*$  and  $\lambda \geq 0$  leads to

$$L(\mathbf{x}^*, \lambda, \mu) = f(\mathbf{x}^*) + \langle \lambda, \mathbf{g}(\mathbf{x}^*) \rangle + \langle \mu, \mathbf{h}(\mathbf{x}^*) \rangle \leq f(\mathbf{x}^*) = L(\mathbf{x}^*, \lambda^*, \mu^*),$$

which completes the proof.

2.  $\Leftarrow$  In view of Theorem 2, it suffices to show that (14) and (17) hold. The left half of the saddle point property of the Lagrangian in (18) implies that

$$\begin{aligned} L(\mathbf{x}^*, \lambda, \mu) &\leq L(\mathbf{x}^*, \lambda^*, \mu^*), \forall \lambda \geq 0, \\ \Rightarrow f(\mathbf{x}^*) + \langle \lambda, \mathbf{g}(\mathbf{x}^*) \rangle + \langle \mu, \mathbf{h}(\mathbf{x}^*) \rangle &\leq L(\mathbf{x}^*, \lambda^*, \mu^*), \forall \lambda \geq 0. \end{aligned}$$

In other words,  $L(\mathbf{x}^*, \lambda, \mu)$  is upper bounded for any  $\lambda \geq 0$ . Consequently, we have

$$\mathbf{g}(\mathbf{x}^*) \leq 0, \mathbf{h}(\mathbf{x}^*) = 0,$$

i.e., the primal feasibility (14) holds (otherwise  $L(\mathbf{x}^*, \lambda, \mu)$  can not be upper bounded).

To show that the complementary slackness in (17) holds, we combine the primal feasibility of  $\mathbf{x}^*$  and left half of (18)

$$\begin{aligned} f(\mathbf{x}^*) + \langle \lambda, \mathbf{g}(\mathbf{x}^*) \rangle &\leq f(\mathbf{x}^*) + \langle \lambda^*, \mathbf{g}(\mathbf{x}^*) \rangle, \forall \lambda \geq 0, \\ \xrightarrow{\lambda \rightarrow 0} \langle \lambda^*, \mathbf{g}(\mathbf{x}^*) \rangle &= \sum_{i=1}^m \lambda_i^* g_i(\mathbf{x}^*) \geq 0. \end{aligned}$$

On the other hand, in view of the facts that  $\lambda^* \geq 0$  and  $\mathbf{g}(\mathbf{x}^*) \leq 0$ , we have

$$\lambda_i^* g_i(\mathbf{x}^*) \leq 0, i = 1, \dots, m.$$

All together, we have

$$\lambda_i^* g_i(\mathbf{x}^*) = 0, i = 1, \dots, m.$$

Thus, the complementary slackness holds and the proof is complete.  $\square$

## 2.5 Strong duality

We discuss conditions that ensure the duality gap is zero.

**Theorem 3.** [1] *Suppose that the primal problem in (1) is a convex optimization problem, that is,  $f$  and  $g_i$ ,  $i = 1, \dots, m$  are convex,  $h_i$ ,  $i = 1, \dots, p$  are affine, and  $X$  is a convex set. If there exists an  $\hat{\mathbf{x}} \in X$  such that  $\mathbf{g}(\hat{\mathbf{x}}) < \mathbf{0}$  and  $\mathbf{h}(\hat{\mathbf{x}}) = \mathbf{0}$ , and  $\mathbf{0} \in \text{int } \mathbf{h}(X)$ , where  $\mathbf{h}(X) = \{\mathbf{h}(\mathbf{x}) : \mathbf{x} \in X\}$ . Then, the duality gap is zero. Furthermore, if  $f^*$  is finite, then there exists at least one geometric multiplier.*

**Proposition 4.** [2] **Strong Duality Theorem - Linear Constraints** *Consider the primal problem. Suppose that  $f$  is convex,  $D_0$  is a polyhedron (that is,  $D_0 = \{\mathbf{x} : \langle \mathbf{a}_i, \mathbf{x} \rangle \leq b_i, i = 1, \dots, r\}$ ), and  $f^*$  is finite. Then, there is no duality gap and there exists at least one geometric multiplier.*

**Proposition 5.** [2] **Linear and Quadratic Programming Duality** *Consider the primal problem. Suppose that  $f$  is convex quadratic,  $D_0$  is a polyhedron, and  $f^*$  is finite. Then, the primal and dual problems have optimal solutions, and the duality gap is 0.*

## 3 The Dual Problem of SVM

### The Primal Problem

Recall that the **soft margin** SVM takes the form of

$$\begin{aligned} \min_{\mathbf{w}, b, \xi} \quad & \frac{1}{2} \|\mathbf{w}\|^2 + C \sum_{i=1}^n \xi_i, \\ \text{s.t.} \quad & y_i(\langle \mathbf{w}, \mathbf{x}_i \rangle + b) \geq 1 - \xi_i, i \in [n], \\ & \xi_i \geq 0, i \in [n]. \end{aligned} \tag{19}$$

The **primal variables** are  $\mathbf{w}$ ,  $b$ , and  $\xi$ . By Proposition (5), the strong duality holds.

### The Lagrangian

To find the dual problem of (19), we first construct the Lagrangian:

$$L(\mathbf{w}, b, \xi, \alpha, \mu) = \frac{1}{2} \|\mathbf{w}\|^2 + C \sum_{i=1}^n \xi_i + \sum_{i=1}^n \alpha_i (1 - \xi_i - y_i(\langle \mathbf{w}, \mathbf{x}_i \rangle + b)) - \sum_{i=1}^n \mu_i \xi_i,$$

where  $\alpha_i, \mu_i \geq 0$ ,  $i = 1, \dots, n$ , are the **dual variables**.



## The Dual Function

We next find the dual function:

$$\begin{aligned}
 q(\alpha, \mu) &= \inf_{\mathbf{w}, b, \xi} L(\mathbf{w}, b, \xi, \alpha, \mu) \\
 &= \inf_{\mathbf{w}} \frac{1}{2} \|\mathbf{w}\|^2 - \sum_{i=1}^n \alpha_i y_i \langle \mathbf{w}, \mathbf{x}_i \rangle \\
 &\quad + \inf_b -b \sum_{i=1}^n \alpha_i y_i \\
 &\quad + \inf_{\xi} \sum_{i=1}^n (C - \alpha_i - \mu_i) \xi_i.
 \end{aligned} \tag{20}$$

For fixed  $(\alpha, \mu)$ , let  $(\hat{\mathbf{w}}, \hat{b}, \hat{\xi})$  be the optimal solution to the above problem. The first order optimal condition implies that

$$\begin{aligned}
 \nabla_{\mathbf{w}} L(\mathbf{w}, b, \xi, \alpha, \mu)|_{\mathbf{w}=\hat{\mathbf{w}}} &= 0 \Rightarrow \hat{\mathbf{w}} - \sum_{i=1}^n \alpha_i y_i \mathbf{x}_i = 0, \\
 \nabla_b L(\mathbf{w}, b, \xi, \alpha, \mu)|_{b=\hat{b}} &= 0 \Rightarrow -\sum_{i=1}^n \alpha_i y_i = 0, \\
 \nabla_{\xi_i} L(\mathbf{w}, b, \xi, \alpha, \mu)|_{\xi_i=\hat{\xi}_i} &= 0 \Rightarrow C - \alpha_i - \mu_i = 0, i = 1, \dots, n.
 \end{aligned}$$

Plugging the above equations into Eq. (20) leads to

$$q(\alpha, \mu) = -\frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j y_i y_j \langle \mathbf{x}_i, \mathbf{x}_j \rangle + \sum_{i=1}^n \alpha_i. \tag{21}$$

## The Dual Problem

Thus, the dual problem of the soft margin SVM in (19) is

$$\begin{aligned}
 \max_{\alpha} \quad & -\frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j y_i y_j \langle \mathbf{x}_i, \mathbf{x}_j \rangle + \sum_{i=1}^n \alpha_i \\
 \text{s.t.} \quad & \sum_{i=1}^n \alpha_i y_i = 0, \\
 & C - \alpha_i - \mu_i = 0, \\
 & \alpha_i \geq 0, \\
 & \mu_i \geq 0, i = 1, \dots, n.
 \end{aligned}$$

We can remove  $\mu$  from the problem by noting that

$$\mu_i = C - \alpha_i, i = 1, \dots, n,$$

which leads to

$$\begin{aligned} \min_{\alpha} \quad & \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j y_i y_j \langle \mathbf{x}_i, \mathbf{x}_j \rangle - \sum_{i=1}^n \alpha_i \\ \text{s.t.} \quad & \sum_{i=1}^n \alpha_i y_i = 0, \\ & \alpha_i \in [0, C], i = 1, \dots, n. \end{aligned} \quad (22)$$

### Complementary Slackness

Let  $(\mathbf{w}^*, b^*, \xi^*)$  and  $(\alpha^*, \mu^*)$  be the optimal solutions to the primal and dual problems of SVM, respectively. By Theorem 2, we write the complementary slackness as follows.

$$\alpha_i^* (1 - \xi_i^* - y_i (\langle \mathbf{w}^*, \mathbf{x}_i \rangle + b^*)) = 0, i = 1, \dots, n, \quad (23)$$

$$\mu_i^* (-\xi_i^*) = (C - \alpha_i^*) (-\xi_i^*) = 0, i = 1, \dots, n. \quad (24)$$

By the complementary slackness in (23) and (24), we have several interesting observations.

1. Suppose that one of the entries of  $\alpha^*$ , say  $\alpha_k^*$ , falls in the interval  $(0, C)$ . Then, the complementary slackness conditions (23) and (24) implies that

$$y_k (\langle \mathbf{w}^*, \mathbf{x}_k \rangle + b^*) = 1 - \xi_k^* \text{ and } \xi_k^* = 0,$$

respectively. Clearly, we have

$$y_k (\langle \mathbf{w}^*, \mathbf{x}_k \rangle + b^*) = 1, \quad (25)$$

which implies that  $\mathbf{x}_k$  is a support vector.

2. Suppose that

$$1 - \xi_k^* - y_k (\langle \mathbf{w}^*, \mathbf{x}_k \rangle + b^*) < 0.$$

Then, by (23) and (24), we have  $\alpha_k^* = 0$  and  $\xi_k^* = 0$ , respectively. Thus,

$$y_k (\langle \mathbf{w}^*, \mathbf{x}_k \rangle + b^*) > 1,$$

which implies that  $\mathbf{x}_k$  is correctly classified and outside of the region between the marginal hyperplanes.

### Recovering the Primal Optimum from the Dual Optimum

**Proposition 6.** *Let  $\alpha^*$  be one of the optimal solutions to (22). Then, we have*

$$\mathbf{w}^* = \sum_{i=1}^n \alpha_i^* y_i \mathbf{x}_i.$$

*If further  $\alpha_k^*$  is one of the entries of  $\alpha^*$  and  $\alpha_k^* \in (0, C)$ , then we have*

$$b^* = y_k - \langle \mathbf{w}^*, \mathbf{x}_k \rangle.$$

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## References

- [1] M. Bazaraa, H. Sherali, and C. Shetty. *Nonlinear Programming*. Wiley-Interscience, 2006.
- [2] D. P. Bertsekas. *Nonlinear Programming, 3ed.* Athena Scientific, 2016.